

ESF Research Conference Future Internet & Society: A Complex Systems Perspective

Interactions among physical, logical and social viewpoints:

An evolutionary design loop

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Complex Networks

- Network: A collection of (nodes, agents, components, objects, services ...) that collaborate to accomplish actions, gains, ...that cannot be accomplished with out such collaboration
- It is all about Interactions that keep increasing and become more complex
- Trade-off: gain from collaboration vs. cost of collaboration
- Complex Networks (CNs): Describe wide range of systems of interacting entities
- Complex Network Analysis
 - Models
 - Properties/features
 - Behavior

Complex Network Taxonomy

Communication, infrastructure, technological networks

Designed and/or engineered

Social and economical networks

Human initiated,
Spontaneous
growth

Biological networks

Spontaneous evolution

Networks: Different Views

- Network Science employs a three level consideration:
 - Physical networks, in which node associations correspond one-to-one in actual interactions among the entities and physical connectivity.
 - Logical networks, involve logical associations and connectivity among peers.
 Such networks include, overlay and peer-to-peer (p2p) networks.
 - Social networks, involves more complex interactions, that take into account mainly unpredictable/hidden social associations (activities).
- Evolutionary Design Loop: Interactions among different

Layers



Control vs. Communications

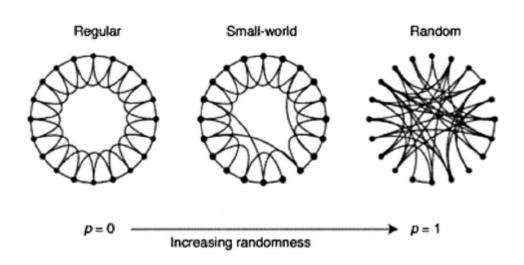
- Many graphs as abstractions
- Collaboration graph or a model of what the system does (behavior)
- Communication graph or a model of what the system consist of (structure)
- Challenge 1: Given behavior, what structure (subject to constraints) gives best performance?
- Challenge 2: Given structure (and constraints) how well behavior can be executed?
- Topology modification topology formation/transformation

Objective

- **Focus**: on closing the loop between social and physical networking in the aforementioned design paradigm.
- Exploit: how social knowledge and features of online social networks can be used in improving the characteristics of physical communication networks
- Demonstrate: infuse the desired properties of an online social structure (small-world effect, power-law like degree distribution) into the core structure of a wireless multi-hop network.
- Use: Inverse Topology Control based techniques to properly add communication links in a multi-hop network.
- Analyze: through a continuum-theory based framework
- Identify: underlying research challenges that need to be addressed for a more holistic treatment and exploitation of the proposed evolutionary design vision.

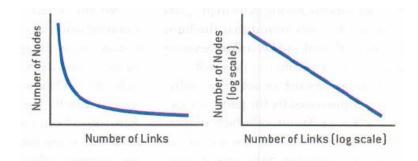
Small-world Networks

- Obtained evolutionary from ordered lattices
 - Start from an ordered lattice
 - Randomly rewire each edge with prob. p excluding selfconnections and duplicate edges
 - Arbitrary long-range edges maybe added
- Small average path length



Scale-free (Exponential) Networks

- Power-law distributed small-world network
- $P(k) \sim k^{-\gamma}$
- Small percentage of nodes with great degrees
- Majority of nodes with small degrees

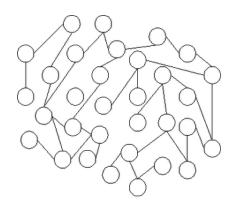


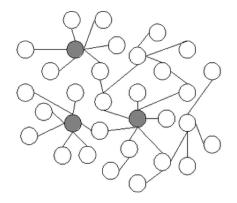
Obtained by growth + preferential attachment

$$\Pi(k_i) = \frac{k_i}{\sum_j k_j}$$

• Many empirically observed networks appear to be scale-free \rightarrow seems the most natural emerging network structure

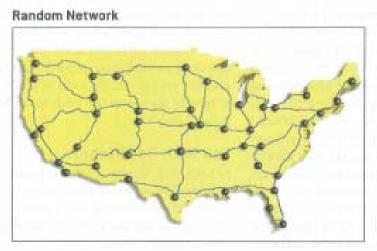
Examples of Scale-free Networks

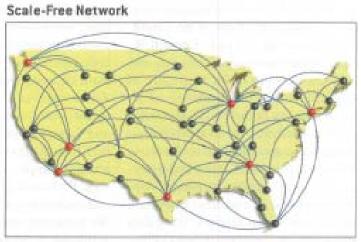




(a) Random network

(b) Scale-free network





Applications of Scale-free Networks

- Internet/WWW
- Science collaboration graphs
- Hollywood co-starring graphs
- Cellular networks
- Citation networks
- road maps
- food chains
- electric power grids
- neural networks
- voter networks
- social influence networks
- Human sexual contacts

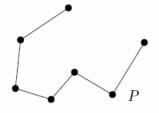
Average Path Length

- Path → sequence of vertices traversed in a network
- Geodesic path ↔ shortest path (topology)
 - shortest path through a network from a vertex to another
- Network diameter → length of longest geodesic

• Definition:
$$l_G = \frac{1}{n*(n-1)} * \sum_{i,j} d(v_i, v_j)$$

- Un-weighted graph
- Total # nodes is n,
- $d(v_i, v_i)$ geodesic length of v_i from v_i





- The actual path length experienced on average by a user
- The lower l_G is, the better it is in general
 - Information dissemination
 - Lower cost

Clustering Coefficient

- Measure of degree → graph nodes tend to cluster together
 - Global

$$C = \frac{3 \times \text{number of triangles}}{\text{number of connected triples of vertices}} = \frac{\text{number of closed triplets}}{\text{number of connected triples of vertices}}.$$

- Local
$$C_i = \frac{|\{e_{jk}\}|}{k_i(k_i-1)} : v_j, v_k \in N_i, e_{jk} \in E.$$

- quantifies how close its neighbors are to being a clique (complete graph)
- (Network-wide) Average Clustering Coefficient
 - Average of the local clustering coefficients of all vertices

$$\bar{C} = \frac{1}{n} \sum_{i=1}^{n} C_i.$$

Inverse Topology Control

• Therefore:

- infuse the scaling behavior of the small-world average path length to a multi-hop network
- exploit <u>Topology Control</u>
- propose various mechanisms for topology modification
- use Continuum Theory for the analysis
- Main objective: basic features of mechanisms to improve selected properties of RGGs making them resemble the behavior of small-world networks (motivated by social network features and processes)
- Advantage in real-time applications: average packet delay (video streaming) and packet loss (QoS)

Description

- At each time step we increase the radius of m_1 selected nodes to a value of R_{max}
- We employ the model at T time steps
- R_{max} , R_{min} : parameters of the network characterizing each time step
- $R_{max}(t+1)=R_{max}(t)+A$
- $R_{min}(t+1)=R_{max}(t)$

Processes

- Process p_1 : With probability p, $0 , we add <math>m_1$, $(m_1 < n)$ new links to selected nodes
- First endpoint: probability $Q_1(k_i)$

Second endpoint: one of the new neighbors in the area of the annulus bounded by R_{min} , R_{max} values of node i

- Probability that a node is selected as the endpoint of the connection: $R(t) = \pi \frac{R_{\text{max}}^2(t) R_{\text{min}}^2(t)}{T^2}$
- **Process** p_2 : With probability (1-p), no change in a node's transmission range $\frac{dk_i}{dt} = pm_1Q_1(k_i)$
- $Q_1(k_i)$: determined by one of the three scenarios

Scenario 1 Preferential attachment to popular nodes

☐ Starting point of the chosen link: one of the nodes with highest node degree (emergence of node-hubs similarly to *scale-free* networks)

$$Q_{1}(k_{i}) = \frac{k_{i} + 1}{\sum_{all_nodes_i} (k_{i} + 1)}$$

- ☐ Preferential attachment regime: nodes favor connections to popular nodes
- ☐ Aim: Reduction of the average path length by connecting nodes of high degree with even more nodes → Better traffic dissemination
- ☐ Concept: based on the scale-free topological nature of most social networks such as the WWW, the networks of scientific paper citations, actors in Hollywood, etc.

Scenario 2

Clustering based on popularity and proximity

☐ Initial endpoint of a new connection chosen regarding: the degree of the node (proportionally), the distance from the center of the deployment area (inversely proportional)

$$Q_{2}(k_{i}) = \frac{k_{i} + 1}{\sum_{all_nodes_i} (k_{i} + 1)} \frac{\frac{1}{d_{i}}}{\sum_{all_nodes_i} \frac{1}{d_{i}}}$$

- ☐ Both popularity and spatial proximity exploited
- ☐ Social analogue: the tension followed by groups of people to cluster on the basis of popularity and geographic proximity

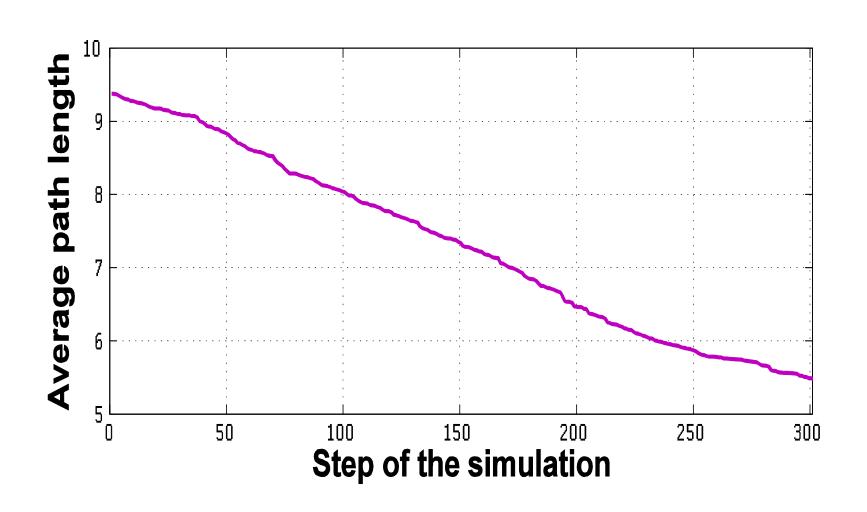
Scenario 3 Preferential attachment with bidirectional links

- Concept similar to Scenario 1 with bidirectional new links added
- Main difference: The transmission radius increases at both the selected node and the nodes to which the initiating nodes connect to
- Bidirectional links lead to higher energy consumption in the whole network but to a larger reduction of the average path length than Scenario 1

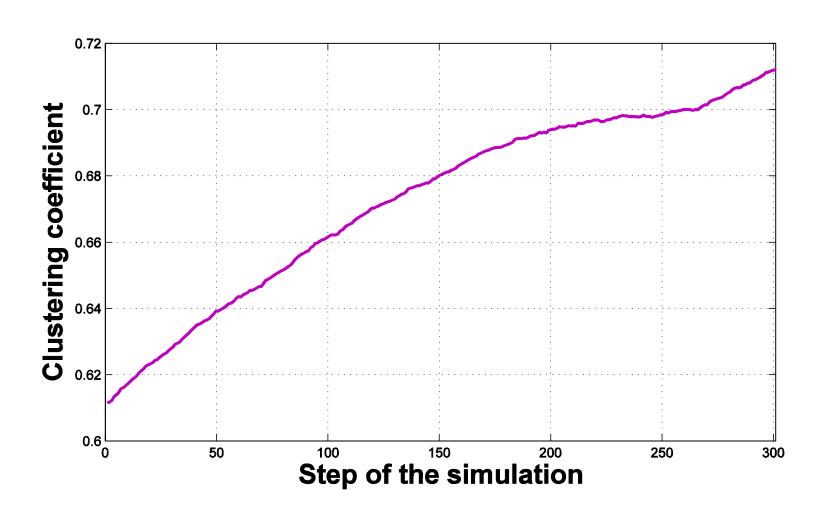
Evaluation – Numerical Results

- Number of nodes N=750.
- Selected nodes at each time step m_1 =10% N=75
- Number of steps=300.
- Radius increment at each time step=1m.
- Square deployment area L=2000m.
- Initial Radius=150m.
- Rmax=150m to 450m.

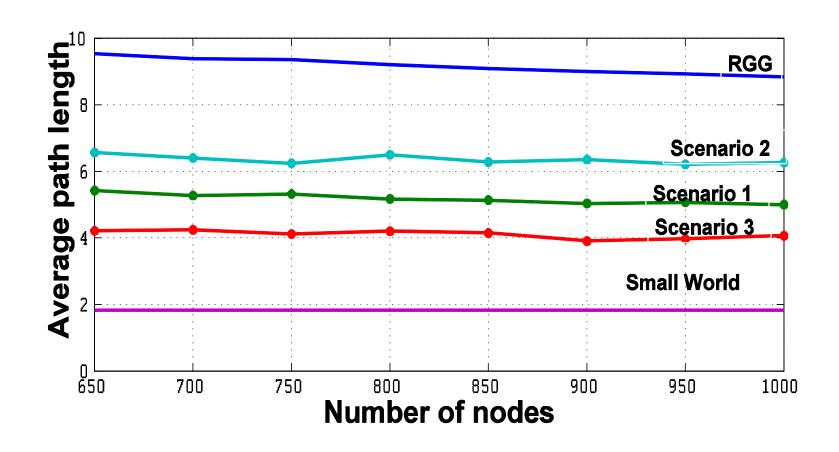
Reduction of Average Path Length



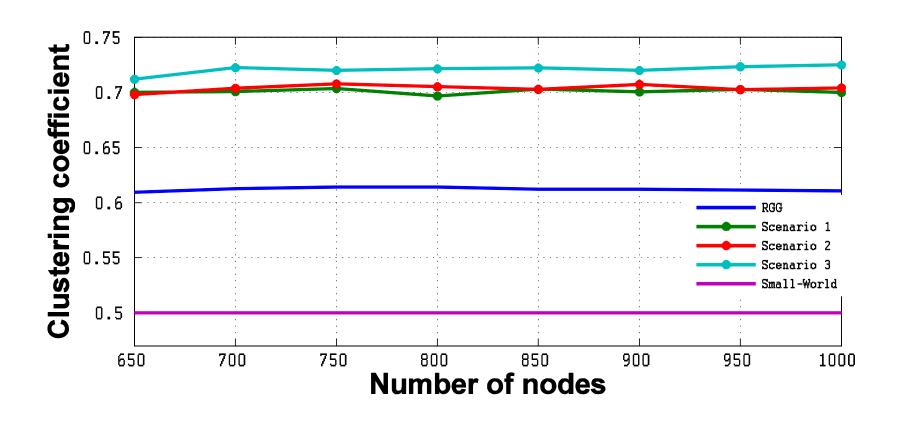
Increase of the Clustering Coefficient



Comparison of Average Path Length



Clustering coefficient comparison



Some recent (2010) relevant publications:

- E. Stai, V. Karyotis and S. Papavassiliou, "Socially-inspired Topology Improvements in Wireless Multi-hop Networks", in Proc. of Social Networks Workshop of IEEE International Conference on Communications (ICC), May 2010.
- E. Stai, V. Karyotis and S. Papavassiliou "Enhanced Service Provisioning in Wireless Multi-hop Networks via Socially-driven Inverse Topology Control", in Proc. of the 4th IEEE Workshop on Enabling the Future Service-Oriented Internet: Towards Socially-Aware Networks, of IEEE Global Communications Conference (GLOBECOM), December 2010.
- E. Stai, V. Karyotis and S. Papavassiliou, "Topology Enhancement in Wireless Multi-Networks: A Top-down Approach", IEEE Trans. on Parallel and Distributed Systems (under review), 2010.
- ➤ V. Karyotis, A. Manolakos and S. Papavassiliou, "On Topology Control and Non-Uniform Node Deployment in Ad Hoc Networks", in Proc. of Sixth IEEE PerCom Workshop on Pervasive Wireless Networking (PWN 2010), April 2010.
- ➤ V. Karyotis and S. Papavassiliou, "Mobility-induced Capacity-Delay Tradeoff in Wireless Multihop Networks", Cluster Computing and Multi-Hop Network Research, Nova Publ., 2010.

Thank you for your attention

Questions?